

ADFOCS 2010, Exercises

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Fisher's market model consists of a set B of n buyers and a set G of g divisible goods. Assume the buyers are numbered from 1 to n and the goods from 1 to g . Let f_i be the function specifying the utility derived by buyer i from a bundle of goods. If x_{ij} is the amount of good j received by buyer i , and \mathbf{x}_i is the g -dimensional vector whose j th component is x_{ij} , then the utility she derives from this bundle is given by $f_i(\mathbf{x}_i)$. Assume buyer i has m_i dollars and w.l.o.g. assume there is 1 unit of each good. The problem is to find *equilibrium prices and allocations*, i.e., prices s.t. if each buyer is given an optimal bundle of goods, the market clears. Several different utility functions will be considered below.

1. **Additively-separable, piecewise-linear, concave utilities:** Assume buyers' utility functions are additively-separable, piecewise-linear, concave, and satisfying non-satiation. Give a polynomial time algorithm for testing if given prices p_1, \dots, p_g are equilibrium prices.

Also, if all parameters are rational and equilibrium exists, show that there must be one that is rational.

2. **Leontief utilities:** A utility function is Leontief if, given parameters $a_{ij} \geq 0$,

$$f_i(\mathbf{x}_i) = \min_j \left\{ \frac{x_{ij}}{a_{ij}} \right\},$$

(assume $0/0$ is large enough to not affect the minimization; alternatively, the minimization is only over non-zero a_{ij} 's). Show that the following convex program captures equilibrium for this utility function. Is this a rational convex program?

$$\begin{aligned} \text{maximize} \quad & \sum_{i \in B} m_i \log(u_i) & (1) \\ \text{subject to} \quad & \forall i \in B, \forall j \in G : u_i \leq \frac{x_{ij}}{a_{ij}} \\ & \forall j \in G : \sum_{i \in B} x_{ij} \leq 1 \\ & \forall i \in B, \forall j \in G : x_{ij} \geq 0 \end{aligned}$$

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3. **Linear utilities:** Show that the following convex program captures Fisher's linear case and that it is rational. Interpret b_{ij} as the amount of money spent by buyer i on good j and p_j as the price of good j .

$$\begin{aligned} \text{maximize} \quad & \sum_{i \in B, j \in G} b_{ij} \log \frac{u_{ij}}{p_j} & (2) \\ \text{subject to} \quad & \forall j \in G : p_j = \sum_{i \in B} b_{ij} \\ & \forall i \in B : \sum_{j \in G} b_{ij} \leq m_i \\ & \forall i \in B, \forall j \in G : b_{ij} \geq 0 \end{aligned}$$